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Stabilizer Slicing: Coherent Error Cancellations in LDPC Codes

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Coherent errors are a dominant noise process in many quantum computing architectures. Unlike stochastic errors, these errors can combine constructively and grow into highly detrimental overrotations. To combat this, we introduce a simple technique for suppressing systematic coherent errors in low-density parity-check (LDPC) stabilizer codes, which we call stabilizer slicing. The essential idea is to slice low-weight stabilizers into two equally-weighted Pauli operators and then apply them by rotating in opposite directions, causing their overrotations to interfere destructively on the logical subspace.

With access to native gates generated by 3-body Hamiltonians, we can completely eliminate purely coherent overrotation errors, and for overrotation noise of 0.99 unitarity we achieve a 135-fold improvement in the logical error rate of Surface-17. For more conventional 2-body ion trap gates, we observe an 89-fold improvement in the logical error rate of Bacon-Shor-13 with purely coherent errors which should be testable in near-term fault-tolerance experiments. This second scheme takes advantage of the prepared gauge degrees of freedom, and to our knowledge is the first example in which the state of the gauge directly affects the robustness of a code’s memory. This work demonstrates that coherent noise is preferable to stochastic noise within certain code and gate implementations when the coherence is utilized effectively.

As we grow closer to experimentally implementing small quantum error-correcting codes [1–3], it becomes increasingly important to study physically motivated error models. While extensive work has been done on simulating the behavior of codes under Pauli noise [4–7], other more realistic models have drawn less attention.

Coherent errors are small unitary operations that occur consistently after every gate, and often require different tools to correct [8–12]. These errors can arise from a variety of sources, and are generally more damaging due to their capacity to grow rapidly when combined [13–16]. In ion trap quantum computing, coherent errors are a dominant noise source, stemming from miscalibrations in experimental equipment, such as laser intensity. This systematic error leads to a relative angle of overrotation, and is a larger source of infidelity than decoherence in most cases.

The intensity of the laser drifts slowly relative to gate times, and so the fraction of overrotation remains approximately constant over a given error correction cycle. This is a double-edged sword: while these errors rapidly accumulate as they are repeatedly applied, they also have a predictable form. Existing proposals take advantage of this uniformity to suppress coherent errors [17]. However these are commonly restricted to improving single-qubit gates due to excessively long gate times when applied to multi-qubit gates [18, 19].

Although we draw our motivation from ion trap quantum computing, the systematic inaccuracy of multi-qubit gates is a major bottleneck for many different architectures. Thus, mitigating coherent errors is an important problem which would benefit any architecture.

Our method takes advantage of experimental degrees of freedom to suppress systematic coherent errors in one of the most important fault-tolerance circuits: syndrome extraction. Importantly, we can reduce the logical error rates without improving the constituent gates or requiring any additional overhead.

To do so, we introduce a technique which we call stabilizer slicing. The essential idea is to split a stabilizer into two equally weighted Pauli rotations, and apply them in opposite directions. In this way, systematic overrotations destructively interfere, leaving only the intended gates.

ARCHITECTURAL REQUIREMENTS

To perform stabilizer slicing, we require a quantum computing architecture with two particular experimental degrees of freedom.

(i) Our architecture gives us the directional freedom to apply any gate in the clockwise or counterclockwise direction.

(ii) For a code with $2n$-body stabilizers, our architecture can generate native multi-qubit gates by evolving an $(n+1)$-body Hamiltonian.

Even when we restricted to standard 2-body interactions (i.e. $n = 1$), our technique has near-term applications to Shor’s code, and to the boundary of surface codes. The directional freedom in (i) can be seen for ion trap multi-qubit Mølmer-Sørenson gates applied to ions of varying interaction parameters, see Figure 1 [20]. In
this case, the freedom of direction can be realized experimentally by adjusting the relative phase of the Raman beams driving the entangling gate. When restricting to low-weight stabilizers in LDPC codes, conditions similar to (ii) are already being proposed for directly implementing multi-qubit measurements in superconducting systems [21]; coupled with subsystem surface code constructions, the reduced circuit volume of syndrome extraction can yield significantly higher thresholds [22].

Given an architecture that satisfies (i) and (ii), we can implement syndrome extraction circuits which are robust to coherent errors. The insight is to use (i) and (ii) to direct corresponding overrotations against each other, thus cancelling them.

STABILIZER SLICING

To illustrate this simple technique, we first consider a leading example. Imagine you are challenged to create the highest fidelity identity channel possible using two applications of a gate $G$ which satisfies $G^2 = I$. Since $G$ is an involution, we can express it as:

$$\exp(-i\theta G) = \cos(\theta)I - i \sin(\theta)G$$  

where $G$ is applied when $\theta = \pm \pi/2$. In the noiseless case, applying two positive rotations is the same as applying one positive and one negative, up to global phase. However, if there exists some coherent overrotation due to miscalibrations in your experimental setup, these cases diverge. In the case where both rotations are in the same direction, errors of the form $\pi/2 \rightarrow (1 + \epsilon)\pi/2$ add constructively, while in the case where the gates are applied in opposite directions, the errors destructively interfere.

While identity circuits don’t come up often, the intuition is similar for our cancellations from stabilizer slicing. When acting on a clean codestate, a stabilizer $S$ effectively acts as an identity operator. Suppose the stabilizer is split into two evenly weighted rotations $S_L$ and $S_R$ satisfying $S_LS_R = S$. Then, if $|\psi\rangle$ is a clean codestate,

$$S|\psi\rangle = |\psi\rangle$$

$$S_R|\psi\rangle = S_L|\psi\rangle$$

$$\therefore \exp(i\theta S_R)|\psi\rangle = \exp(i\theta S_L)|\psi\rangle.$$  

Following a similar intuition as the previous case, we can apply our stabilizer through two controlled $\pi/2$-rotations. If our errors are overrotations of the form

$$U^L_E := \exp(i\epsilon \theta_L CS_L)$$

$$U^R_E := \exp(i\epsilon \theta_R CS_R),$$  

then by Equation 2 we can replace $S_R$ with $S_L$ and have $\theta_L$ and $\theta_R$ point in opposite directions in order to completely cancel our errors.

$$U^L_E U^R_E |\psi\rangle + \rangle = \exp(i\epsilon \theta_L CS_L) \exp(i\epsilon \theta_R CS_R) |\psi\rangle + \rangle$$

$$\exp(i\epsilon \theta_L CS_L) \exp(i\epsilon (-\theta_L)CS_L) |\psi\rangle + \rangle$$

$$= |\psi\rangle + \rangle.$$  

This technique of applying a stabilizer in two halves which have perfectly cancelling errors is what we call stabilizer slicing, and the corresponding circuit is shown in Figure 2. Of course, we perform syndrome extraction to detect errors, and so sometimes $|\psi\rangle$ will not be a clean codestate. If some stochastic error $E$ occurs on our data, it may put the data into a state for which $S|\psi\rangle = -|\psi\rangle$. In this case, a stabilizer sliced circuit would actually grow coherent errors when measuring violated stabilizers. In this way, we can see that coherent error suppression via stabilizer slicing interfaces non-trivially with other sources of noise. However, in the low-error regime, the majority of stabilizers will commute with $E$ and so stabilizer slicing will have an overwhelmingly positive effect on coherent errors in total.

It is worth noting why property (ii) is necessary. Although we can express any such controlled-$S$ as a product of noiseless two-qubit gates, overrotations on the individual gates will not perfectly cancel. For a weight-$n$ stabilizer $S$ expanded as a product of $m \geq 3$ multi-qubit rotations, interference only occurs between weight $k$ and weight $n-k$ components. Near $\theta = \pm \pi/2$, this yields only a negligible suppression of the two-qubit coherent errors. Thus, splitting the stabilizer into precisely two native components is essential for perfect cancellation.
To best display the improvements of stabilizer slicing, we cal error rates due to the dropout of off-diagonal terms. We reiterate that, ordinarily, we would expect the latter channel, corresponding to $\kappa$, to produce lower logical error rates than the former channel, corresponding to $\epsilon$. The first parameter is the unitarity $\kappa$, and the second parameter is the overrotation angle $\epsilon$, which controls the strength of the error. Consequently, the error following some perfect gate $G$ has the form,

$$\varepsilon_G(\rho) = \kappa \cdot \varepsilon_G^\epsilon(\rho) + (1 - \kappa) \cdot \varepsilon_G^\kappa(\rho)$$

where $\varepsilon_G^\epsilon$ and $\varepsilon_G^\kappa$ are coherent and stochastic overrotation channels with equal fidelity given by,

$$\varepsilon_G^\epsilon(\rho) = \exp(-iG)\rho\exp(iG)$$
$$\varepsilon_G^\kappa(\rho) = \cos^2(\epsilon)I\rho I + \sin^2(\epsilon)G\rho G.$$ (6)

We reiterate that, ordinarily, we would expect the latter channel, corresponding to $\kappa = 0$, to produce lower logical error rates due to the dropout of off-diagonal terms. To best display the improvements of stabilizer slicing, we do not include measurement error. Since coherent error is difficult to simulate, we must use more computationally intensive techniques. Using the full density simulator quantumsim, we write our error correcting circuits as a full quantum channel [3]. By postselecting on all possible syndrome outcomes and then combining the resulting sub-normalized density matrices, we are able to calculate exact logical error rates. While this technique is perfectly accurate, it is exponential in the number of postselections, and consequently we are restricted to simulations with perfect preparation of logical states.

The first simulation we show is for a stablizer sliced Surface-17 code. In this simulation we assume access to 3-body $CXX$ and $CZZ$ gates, and have an error model where $\epsilon_2 = \epsilon_3$ for simplicity. It should be noted that since the 2-body and 3-body gates are never applied within the same stabilizer, this condition is not important for the effectiveness of our scheme. As can be seen in Figure 3, interaction between the stochastic and coherent errors leads to a non-linear interpolation as we move from stochastic to coherent.

Our next simulation for Bacon-Shor assumes current technologies and implements a more physically grounded error model and gate set. We use the decompositions in Figure 1 to convert our 6-body Bacon-Shor stabilizers into ion trap gates, as in Figure 5. In this model, we consider overrotations on both one- and two-qubit gates, with $1 + \epsilon_2 = (1 + \epsilon_1)^2$ reflecting the quadratic dependence of the two-qubit Rabi frequency on the one-qubit Rabi frequency [34].

In order to make this a more feasible system, we restrict our qubit number and consequently can only measure large stabilizers in parallel. As a result, instead of slicing stabilizers, we are limited to slicing gauges instead.
This is an issue as our cancellations rely on the eigenvalue of the operator we are slicing to be +1. As a result we can prepare into a $|0\rangle_L$ state with all $X$-gauges being equal to +1, but over time, $Z$-errors and corrections will flip these $X$-gauges and lead to a degradation of our cancellations. To understand the impact that this gauge decay has on our system, we consider multiple rounds of error correction in Figure 6. By occupying an $X$-type gauge we occupy a superposition over all $Z$-type gauge eigenstates. Consequently, there will be no suppression of $Z$-type coherent errors. As a result the error shown in the Figures 4 is the one-sided infidelity of $|0\rangle_L$.

**CONCLUSIONS**

In both simulations, stabilizer slicing shows marked improvements in the logical error rate, and could be extended to several other codes without modification. We would like to point out a few cases where stabilizer slicing applies in a wider context.

First, note that our decomposition into two components needn’t be symmetric. Although we have described stabilizer splitting as occurring between two equally-weighted (and in Figure 2, disjoint) sets of Pauli operators, the only requirement is that $S = S_L S_R$.

However, to remain experimentally motivated, we wouldn’t expect that systematic overrotations between different many-body interactions would be of the same approximate magnitude. While this is physically realistic for the same process mediating the same multi-qubit interaction among different subsets of qubits, processes mediating different weight interactions will likely have different relative overrotations leading to imperfect cancellations.

It is worth noting that we can still realize evenly-weighted decompositions of odd-weight stabilizers. One simple example is a twist defect, which enables efficient logical Clifford operations on the surface code. A twist defect is formed by odd-weight stabilizers $XXX$. With access to native $CXX$ and $CZZ$ gates, we can slice such a stabilizer as $S_L = XXX$ and $S_R = IZZZ$ and reasonably expect coherent error cancellation.

Stabilizer slicing also extends to Shor-style syndrome extraction using a Bell state $|\Phi^+\rangle$. This follows from the Bell state itself satisfying $ZZ |\Phi^+\rangle = XX |\Phi^+\rangle = |\Phi^+\rangle$. By using native multi-qubit interactions, we are already implicitly saving on the circuit-depth of syndrome extraction. Using such a scheme would allow syndrome-extraction in a single gate-layer.

The downside of this approach is that, even using Bell states, we may introduce multi-qubit correlated errors due to failures in native multi-qubit gates. However, for LDPC codes on which these few-body stabilizer interactions may be reasonable, these correlated errors cannot propagate too badly. In certain cases, such as the hook...
errors of the surface code, such errors will not lower the effective distance of the code. Furthermore, for architectures where stochastic depolarizing errors due to gate failure are rare, coherent error mitigation may be well worth the trade-off.

Lastly, we note that the concept of stabilizer slicing can be applied to any circuit suffering from systematic coherent errors, with varying efficacy. In particular, preparation circuits with a fixed input may be a good candidate to extend stabilizer slicing beyond syndrome extraction.

In summary, stabilizer slicing is a new and simple technique for suppressing coherent errors in syndrome extraction. It requires certain experimental capacities but no additional overhead, and dramatically improves the logical fidelity of syndrome extraction with the same-quality physical gates. Because it requires no additional resources, we hope that even its 2-body iteration could yield significant benefit in realistic near-term fault-tolerance experiments where systematic coherent error is a dominant factor.

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